



Considering Network Heterogeneity in Global Application Layer Multicast Provision

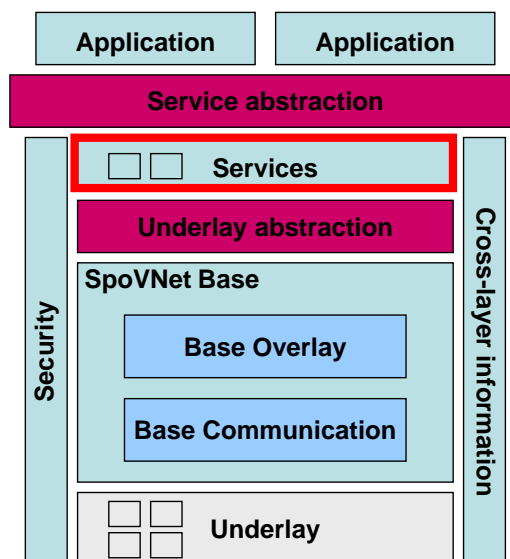
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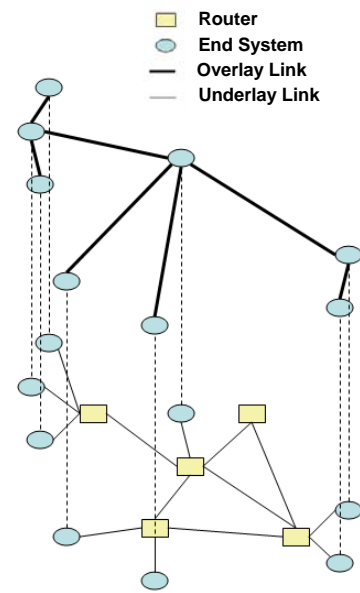
Getting higher in the SpoVNet stack...



- Services in SpoVNet
 - Reside above Underlay Abstraction
- Offer interface to applications
- Several services as part of the architecture
 - MCP-O (Data Dissemination)
 - ES (Event Service)
 - Security Service
- Focus here: MCP-O

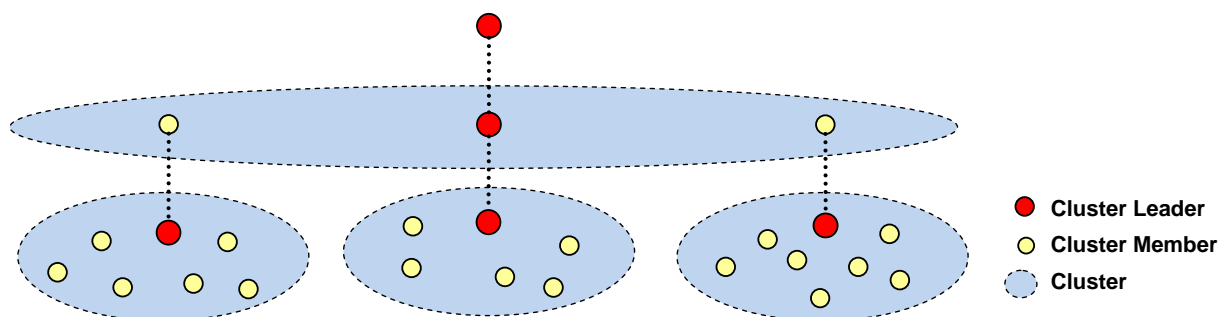


- **Group Communication** in SpoVNet instances
 - Solution: Application Layer Multicast (ALM)
- Several existing approaches
 - Narada, NICE, Yoid, Nemo, ...
 - Each with different target applications
- Drawbacks
 - Considering **homogeneous** (UDP-)underlays
 - Measure link property themselves
 - Mostly latency
- **SpoVNet Vision**
 - Global service provision
 - **Heterogeneity** of nodes and network!
 - Different **application demands**

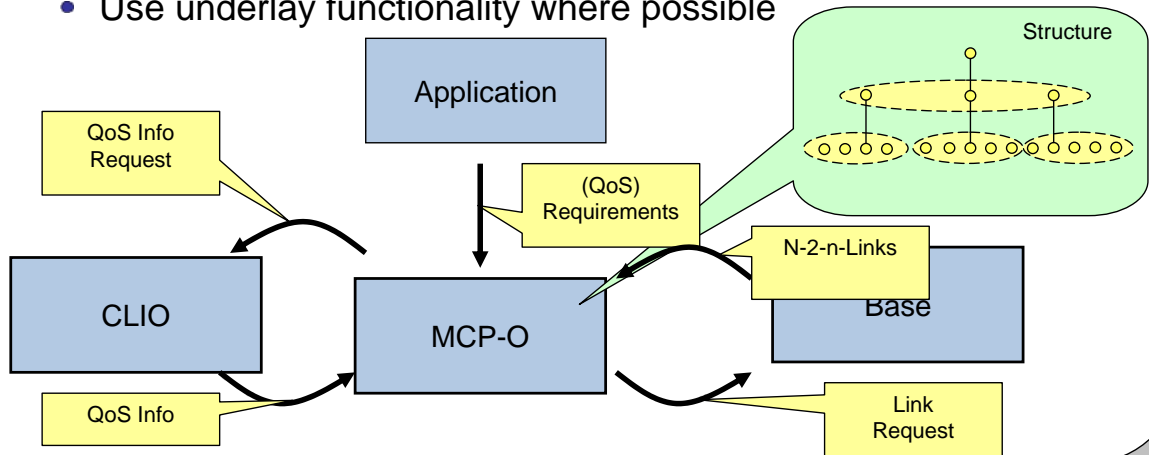


Application Layer Multicast (ALM)

- Achieve scalability through **hierarchical clustering**
 - Limit protocol overhead
- Consider node distances (n-2-n latencies)
- Logical structure determines data dissemination



- Basic idea (enhance NICE approach to...)
 - Evaluate application demands for structural design decisions
 - Consider Cross-Layer-Information (provided by CLIO comp.)
 - Built efficient dissemination topologies for every use case
 - Use underlay functionality where possible

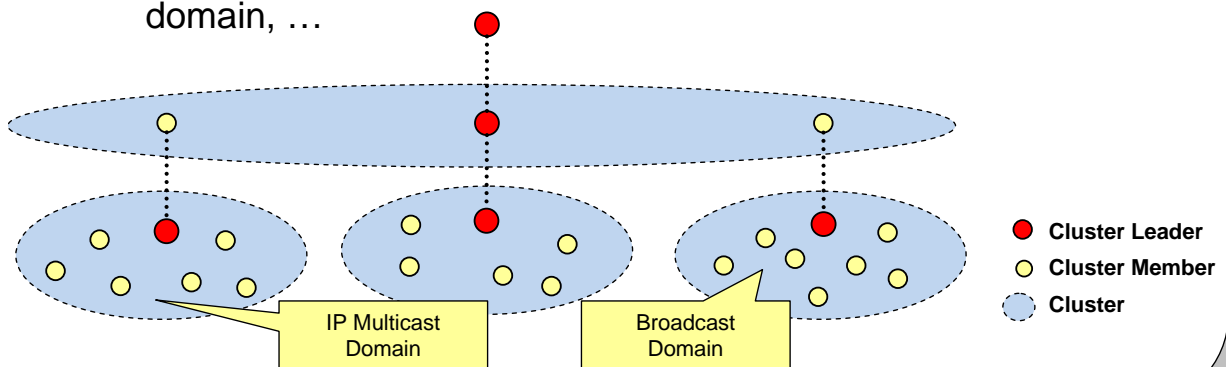


- Use **Service Metric** to affect building of structure

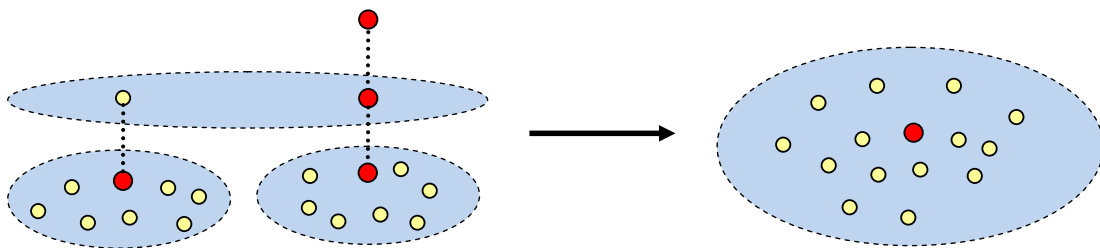
$$d : K \times K \rightarrow \mathbb{R}_{\geq 0}$$

$$d : (x, y) \mapsto \sum_i a_i f_i(x, y)$$

- Weights a_i determined through application requirements
- $f_i(x,y)$ e.g. latency, same WiFi domain, same IP Multicast domain, ...

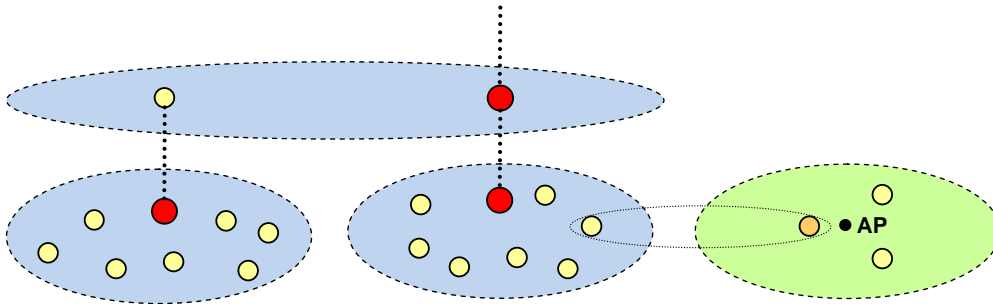


- But: Hierarchical approach not always the best choice
 - Good scalability, but cutting back e.g. latencies ☺
- Example: Latency-sensitive app with < 16 Members
- Adapt cluster size K
 - $K = (\text{lowest current upstream}) / (\text{data stream})$
 - Unicast brings best achievable latency, protocol decides reasonable overhead in bounds

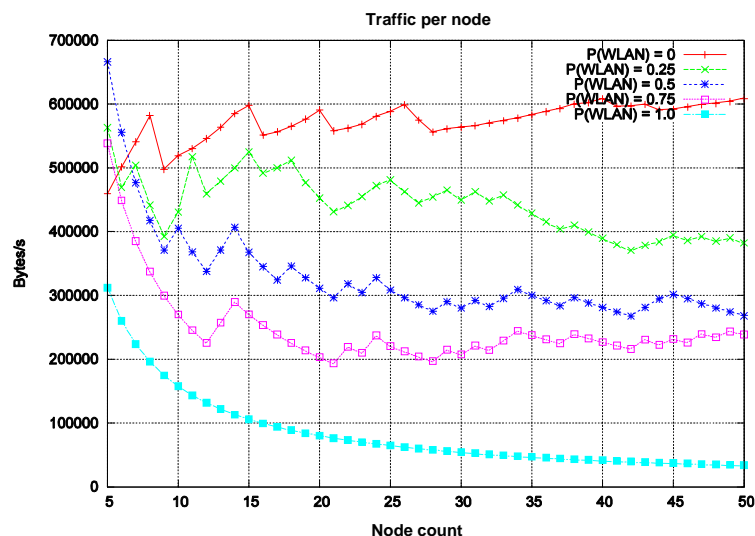


- Single Source Videostreaming application
 - Locate High-Bandwidth nodes near source
 - Make them cluster leader, upstream determines cluster size
 - Flat tree to bound latency and jitter
- Realtime Game
 - Latency critical factor
 - Use unicast where feasible

- First Steps: Wireless Integration
- Save transmission overhead through broadcast
→ NICE-WLI
- Introduce **Gateway-Nodes** for mediation
 - WiFi domain not part of structure, but connected through Gateway
 - Prevent Gateways from becoming cluster leaders (save overhead)



- Approach saves data transmission overhead
 - Figure: $P(\text{WLAN})$ denotes probability of WiFi node



- Efficient data transmission for SpoVNet applications through
 - Considering node and network heterogeneity
 - Use appropriate dissemination strategies
 - IP Multicast, broadcast, QoS reservations,...
- WiFi support already integrated

- Further Work
 - Multiple group support per instance
 - Usage of underlay features